# Mathematics 3670: Computer Systems Bits, Data Types, and Operations 

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| What | When |
| :--- | :--- |
| Read Chapter 2 | this week |
| Design Lab 2 TMs | before Thursday |
| Complete Lab 2 | this Thursday |
| Submit Lab 2 work | by next Thursday |

## Geek humor



Source: http://xkcd.com/571/


3 bits yield $2^{3}=8$ possibilities

## Number of bit patterns

| number of bits | number of bit patterns |
| :---: | :---: |
| 3 | $2^{3}=8$ |
| 4 | $2^{4}=16$ |
| $\ldots$ | $\ldots$ |
| $m$ | $2^{m}$ |
| $\ldots$ | $\ldots$ |
| 16 | $2^{16}=65,536$ |
| $\ldots$ | $\ldots$ |
| 32 | $2^{32}=4,294,967,296$ |
| $\ldots$ | $\ldots$ |
| 64 | $2^{64}=18,446,744,073,709,551,616$ |
| $\ldots$ | $\ldots$ |

## Representations

What is the meaning of 0011010111110010 ?

- An integer? If so, which representation?
- One or more characters? (ASCII or Unicode)
- A floating point value?
- A value of an enumeration type?
- Something else?

| Consider a bit string such as: 0011010111110010 |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Use 4-bit groups |  |  | 0011010111110010 |  |  |  |  |
| Use hexadecimal digits: |  |  | 35 F 2 |  |  |  |  |
| pattern | 0000-1001 | 1010 | 1011 | 1100 | 1101 | 1110 | 1111 |
| hexadecimal | 0-9 | A | B | C | D | E | F |

## ASCII code

- American Standard Code for Information Interchange
- ASCII uses a 7-bit code
- 7 bits allows for only $2^{7}=128$ different characters
- See http://highered.mcgraw-hill.com/sites/dl/free/ 0072467509/104653/PattPatelAppE.pdf


## Unicode

- One system for all the world's languages
- Unicode uses a muti-byte code
- 2 bytes provides $2^{16}=65,536$ different characters
- See http://www.unicode.org/charts/


## Wheel of 3-bit codes: food choices (enumeration type)



## Wheel of 3-bit codes: unsigned integers



## Wheel of 3-bit codes: signed magnitude integers



## Wheel of 3-bit codes: one's complement integers



## Wheel of 3-bit codes: two's complement, signed integers



## Wheel of $m$-bit codes: two's complement, signed integers



## Lab 2 exercise: complement and add one

$$
\begin{array}{cllllllllll}
n & \text { input } & 0 & 1 & 1 & 0 & 0 & 1 & 0 & 0 & 0 \\
& & & & & & \Downarrow & & & & \\
& \text { complement } & 1 & 0 & 0 & 1 & 1 & 0 & 1 & 1 & 1 \\
& & & & & & \Downarrow & & & & \\
& \text { add one } & 1 & 0 & 0 & 1 & 1 & 1 & 0 & 0 & 0
\end{array}
$$

Consider $n+T C(n) \ldots$

$$
\begin{array}{llllllllll}
n & 0 & 1 & 1 & 0 & 0 & 1 & 0 & 0 & 0 \\
T C(n) & 1 & 0 & 0 & 1 & 1 & 1 & 0 & 0 & 0 \\
\hline & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0
\end{array}
$$

$T C(n)$ is the additive inverse of $n$ : i.e., $n+T C(n)=0$

## Two's complement addition

Let $m=b_{n-1} b_{n-2} \ldots b_{2} b_{1} b_{0}$ be an arbitrary $n$-bit pattern
Complement each bit: $C(m)=\bar{b}_{n-1} \bar{b}_{n-2} \ldots \bar{b}_{2} \bar{b}_{1} \bar{b}_{0}$

$$
\begin{aligned}
m+T C(m) & =m+(C(m)+1) \\
& =(m+C(m))+1 \\
& =\left(b_{n-1} b_{n-2} \ldots b_{2} b_{1} b_{0}+\bar{b}_{n-1} \bar{b}_{n-2} \ldots \bar{b}_{2} \bar{b}_{1} \bar{b}_{0}\right)+1 \\
& =(11 \ldots 111)+1 \\
& =00 \ldots 000
\end{aligned}
$$

## Conclusion

If $m$ represents an integer $k$, then $T C(m)$ represents $-k$.

Two's complement: example

Using an 8-bit register, what is the two's complement representation of -20 ?

$$
\left.\begin{array}{rl}
20 & =16+4 \\
& =0
\end{array}\right)
$$

Verify...

$$
\begin{array}{lllllllll}
n & 0 & 0 & 0 & 1 & 0 & 1 & 0 & 0 \\
T C(n) & 1 & 1 & 1 & 0 & 1 & 1 & 0 & 0 \\
\hline
\end{array}
$$

Two's complement: example

Using an 8-bit register, what is $T C(-20)$ ?

$$
\left.\begin{array}{rl}
-20 & \rightarrow \\
\text { complement } & \rightarrow \\
\text { add one } & \rightarrow 0 \\
& =0
\end{array}\right)
$$

Given a bit string $n=b_{w-1} b_{w-2} \ldots b_{2} b_{1} b_{0}$, what value is represented?

| MSB? | Conclusion | What to do |
| :--- | :--- | :--- |
| $b_{w-1}=0$ | value is non-negative | evaluate $n$ as a binary value |
| $b_{w-1}=1$ | value is negative | find $T C(n)$, evaluate, affix sign |$\quad$| 0 | 1 | 0 | 1 | 0 | $0=?$ |  |  |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 1 | 0 | 0 | 1 | 1 | 0 | 0 | $0=?$ |

## Two's complement: decimal to binary conversion

Give a decimal value $n$ and a word length $w$, what bit string $b_{w-1} b_{w-2} \ldots b_{2} b_{1} b_{0}$ represents $n$ ?

| Sign? | What to do |
| :--- | :--- |
| $n \geq 0$ | convert $(n)$ |
| $n<0$ | $\operatorname{TC}(\operatorname{convert}(\|n\|))$ |

To convert a non-negative value...

Greedy "brute force" algorithm identify highest powers of two Successive division by two identify bits from LSB to MSB

Examples: Using an 8-bit word...

$$
\begin{aligned}
57 & =? \\
-57 & =?
\end{aligned}
$$

## Decimal to binary conversion (non-negative value)

## convert $(n)$

Successive division by two generates the bits in reverse order, from LSB to MSB. For example, take $n=57$ :

$$
\begin{aligned}
57 \div 2 & =28 \times 2+1 \\
28 \div 2 & =14 \times 2+0 \\
14 \div 2 & =7 \times 2+0 \\
7 \div 2 & =3 \times 2+1 \\
3 \div 2 & =1 \times 2+1 \\
1 \div 2 & =0 \times 2+1
\end{aligned}
$$

Conclusion: $(57)_{10}=(111001)_{2}$. For an 8 -bit register, we fill with leading zeros: $(57)_{10}=(00111001)_{2}$.

## Decimal to binary conversion (negative value)

TC(convert (|n|))
What is the 8 -bit, two's complement representation of $n=-57$ ?

$$
\begin{aligned}
\operatorname{convert}(|n|) & =\operatorname{convert}(57) \\
& =(00111001)_{2}
\end{aligned}
$$

Now, find the two's complement. . .

$$
\begin{array}{rlllllllll}
57 & \rightarrow & 0 & 0 & 1 & 1 & 1 & 0 & 0 & 1 \\
\text { complement } & \rightarrow & 1 & 1 & 0 & 0 & 0 & 1 & 1 & 0 \\
\text { add one } & \rightarrow & 1 & 1 & 0 & 0 & 0 & 1 & 1 & 1
\end{array}
$$

## Addition on binary quantities

| $a_{3}$ | $a_{2}$ | $a_{1}$ | $a_{0}$ |
| :---: | :---: | :---: | :---: |
| $b_{3}$ | $b_{2}$ | $b_{1}$ | $b_{0}$ |

## Addition on binary quantities

|  | $c_{1}$ |  |  |
| :---: | :---: | :---: | :---: |
| $a_{3}$ | $a_{2}$ | $a_{1}$ | $a_{0}$ |
| $b_{3}$ | $b_{2}$ | $b_{1}$ | $b_{0}$ |
|  |  |  | $s_{0}$ |

## Addition on binary quantities

|  | $c_{2}$ | $c_{1}$ |  |
| :---: | :---: | :---: | :---: |
| $a_{3}$ | $a_{2}$ | $a_{1}$ | $a_{0}$ |
| $b_{3}$ | $b_{2}$ | $b_{1}$ | $b_{0}$ |
|  |  | $s_{1}$ | $s_{0}$ |

## Addition on binary quantities

| $c_{3}$ | $c_{2}$ | $c_{1}$ |  |
| :---: | :---: | :---: | :---: |
| $a_{3}$ | $a_{2}$ | $a_{1}$ | $a_{0}$ |
| $b_{3}$ | $b_{2}$ | $b_{1}$ | $b_{0}$ |
|  | $s_{2}$ | $s_{1}$ | $s_{0}$ |

## Addition on binary quantities

$c_{4}$| $c_{3}$ | $c_{2}$ | $c_{1}$ |  |
| :---: | :---: | :---: | :---: |
| $a_{3}$ | $a_{2}$ | $a_{1}$ | $a_{0}$ |
| $b_{3}$ | $b_{2}$ | $b_{1}$ | $b_{0}$ |
| $s_{3}$ | $s_{2}$ | $s_{1}$ | $s_{0}$ |

## Addition on binary quantities

$c_{4}$| $c_{3}$ | $c_{2}$ | $c_{1}$ |  |
| :---: | :---: | :---: | :---: |
|  | $a_{3}$ | $a_{2}$ | $a_{1}$ |
|  | $a_{0}$ |  |  |
| $b_{3}$ | $b_{2}$ | $b_{1}$ | $b_{0}$ |
|  | $s_{3}$ | $s_{2}$ | $s_{1}$ |
|  | $s_{0}$ |  |  |

We ignore the "carry out" $c_{4}$ generated in the leftmost column

## Addition on binary quantities: example 1

| 0 | 0 | 1 | 1 |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 1 | 0 |

## Addition on binary quantities: example 1



## Addition on binary quantities: example 1



## Addition on binary quantities: example 1

| 0 | 1 | 0 |  |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 1 | 1 |
| 0 | 0 | 1 | 0 |
|  | 1 | 0 | 1 |

## Addition on binary quantities: example 1

| 0 | 0 | 1 | 0 |  |
| :--- | :--- | :--- | :--- | :--- |
|  | 0 | 0 | 1 | 1 |
|  | 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 1 |  |

## Addition on binary quantities: example 1

| 0 | 0 | 1 | 0 |  |
| :--- | :--- | :--- | :--- | :--- |
|  | 0 | 0 | 1 | 1 |
|  | 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 1 |  |

This shows $3+2=5$ in a 4-bit system

## Addition on binary quantities: example 2

| 0 | 0 | 1 | 1 |
| :--- | :--- | :--- | :--- |
| 1 | 0 | 1 | 1 |

## Addition on binary quantities: example 2



## Addition on binary quantities: example 2

|  | 1 | 1 |  |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 1 | 1 |
| 1 | 0 | 1 | 1 |
|  |  | 1 | 0 |

## Addition on binary quantities: example 2

| 0 | 1 | 1 |  |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 1 | 1 |
| 1 | 0 | 1 | 1 |
|  | 1 | 1 | 0 |

## Addition on binary quantities: example 2

| 0 | 0 | 1 | 1 |  |
| :--- | :--- | :--- | :--- | :--- |
|  | 0 | 0 | 1 | 1 |
|  | 1 | 0 | 1 | 1 |
|  | 1 | 1 | 1 | 0 |

## Addition on binary quantities: example 2

$$
\begin{array}{lllll}
0 & 0 & 1 & 1 & \\
& 0 & 0 & 1 & 1 \\
& 1 & 0 & 1 & 1 \\
\hline & 1 & 1 & 1 & 0
\end{array}
$$

This shows $3+(-5)=-2$ in a 4-bit system

## Addition on binary quantities: example 3

| 0 | 1 | 0 | 1 |
| :--- | :--- | :--- | :--- |
| 0 | 1 | 0 | 0 |

## Addition on binary quantities: example 3



## Addition on binary quantities: example 3

|  | 0 | 0 |  |
| :--- | :--- | :--- | :--- |
| 0 | 1 | 0 | 1 |
| 0 | 1 | 0 | 0 |
|  |  | 0 | 1 |

## Addition on binary quantities: example 3

| 1 | 0 | 0 |  |
| :--- | :--- | :--- | :--- |
| 0 | 1 | 0 | 1 |
| 0 | 1 | 0 | 0 |
|  | 0 | 0 | 1 |

## Addition on binary quantities: example 3

| 0 | 1 | 0 | 0 |  |
| :--- | :--- | :--- | :--- | :--- |
|  | 0 | 1 | 0 | 1 |
|  | 0 | 1 | 0 | 0 |
| 1 | 0 | 0 | 1 |  |

## Addition on binary quantities: example 3

0 | 1 | 0 | 0 |  |  |
| :--- | :--- | :--- | :--- | :--- |
|  | 0 | 1 | 0 | 1 |
|  | 0 | 1 | 0 | 0 |
| 1 | 0 | 0 | 1 |  |

This shows $5+4=-7$ in a 4 -bit system Oops: arithmetic overflow

## Overflow summary for $A+B$

| $A$ | $B$ | Outcome |
| :---: | :---: | :--- |
| positive | negative | correct result |
| negative | positive | correct result |
| negative | negative | possible overflow |
| positive | positive | possible overflow |

Informal justification: two's complement wheel

## Bit fiddling: arithmetic left shift

Various low-level operations on bit strings are often useful

## Arithmetic left shift

$$
b_{7} b_{6} b_{5} b_{4} b_{3} b_{2} b_{1} b_{0} \text { becomes } b_{6} b_{5} b_{4} b_{3} b_{2} b_{1} b_{0} 0
$$

If there is no overflow. . .

- an arithmetic left shift operation computes $2 k$, given $k$
- $n$ successive arithmetic left shifts computes $2^{n} k$, given $k$


## Bit fiddling: sign extension

- We use sign extension when we increase the number of bits
- For example, we may convert an 4-bit value to a 8 -bit value
- Simply replicate the MSB
- $b_{3} b_{2} b_{1} b_{0}$ becomes $b_{3} b_{3} b_{3} b_{3} b_{3} b_{2} b_{1} b_{0}$

$$
\begin{array}{lll}
0101 & 00000101 & +5 \\
1101 & 11111101 & -3
\end{array}
$$

Why does it work?

## Bit fiddling: bitwise AND

| $a$ | $b$ | $a$ AND $b$ |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 0 |
| 1 | 0 | 0 |
| 1 | 1 | 1 |

Summary

- 1 AND $b=b$
- 0 AND $b=0$


## Bit fiddling: bitwise OR — inclusive or

| $a$ | $b$ | $a$ OR $b$ |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 1 |
| 1 | 0 | 1 |
| 1 | 1 | 1 |

Summary

- 1 OR $b=1$
- 0 OR $b=b$


## Bit fiddling: bitwise XOR — exclusive or

| $a$ | $b$ | $a$ XOR $b$ |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 1 |
| 1 | 0 | 1 |
| 1 | 1 | 0 |

Summary

- $a=b$ yields 0
- $a \neq b$ yields 1


## Bit fiddling: masking operations

AND is useful for isolating specific bits

\[

\]

OR is useful for inserting ones

$$
\begin{array}{ccccccccc} 
& b_{7} & b_{6} & b_{5} & b_{4} & b_{3} & b_{2} & b_{1} & b_{0} \\
\text { OR } & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 0 \\
\cline { 2 - 8 } & 1 & 1 & 1 & 1 & 1 & b_{2} & b_{1} & b_{0}
\end{array}
$$

## Bit fiddling: XOR application - testing for equality



Two bit patterns match if and only if all result bits are 0

## Floating point representation

- Use a fixed number of bits, e.g., 32 bits
- Subdivide bits into fields: sign, exponent, fraction
- IEEE floating point standard (including Java's float):
- 1 sign bit
- 8 exponent bits, using "excess 127 "
- 23 fraction bits plus "hidden bit"
- Example 1: How can we represent $-6 \frac{5}{8}$ as a 32-bit float?
- Example 2: What float value is represented by 3D800000 ?

